

Performance of randomized forwarding methods in large ad hoc networks

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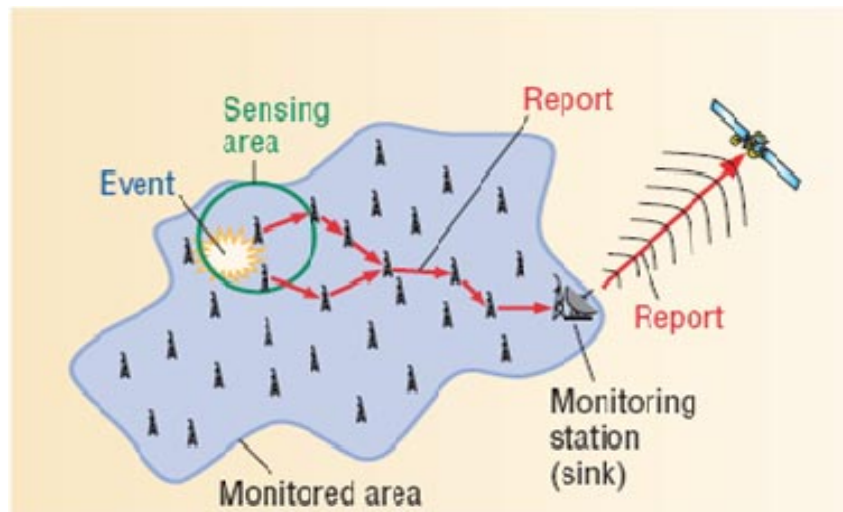
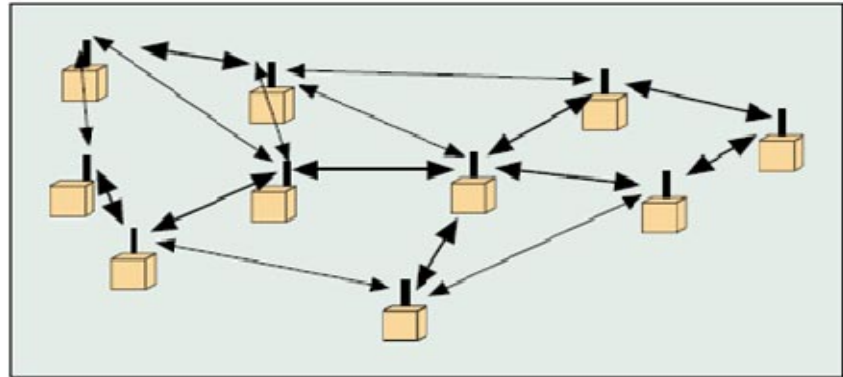
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CONTENTS

- Ad hoc networks
- Objective of the study
- MAC and routing in ad hoc networks
- Network model for simulations
- Used forwarding algorithms
- Results
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AD HOC NETWORKS

- Wireless nodes that communicate without fixed infrastructure or centralized control
- Multihop communications
 - Each node acts as a router
- Military and rescue applications
- Wireless sensor networks
 - Sensing, data processing and communications



OBJECTIVE OF THE STUDY

- Survey of MAC and routing methods in ad hoc networks
 - Clear classification and most important methods
- Simulation study to maximize network-wide throughput and to compare the performance of geographic forwarding methods in a large ad hoc network
 - Maximization of packet flow intensity with respect to network density and the slotted ALOHA transmission probability

MEDIUM ACCESS CONTROL

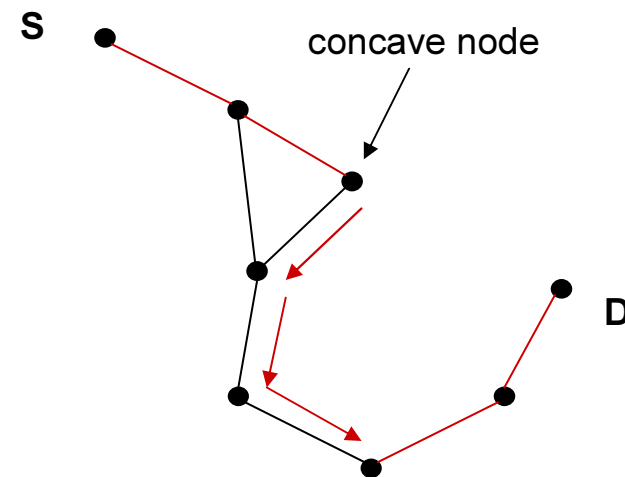
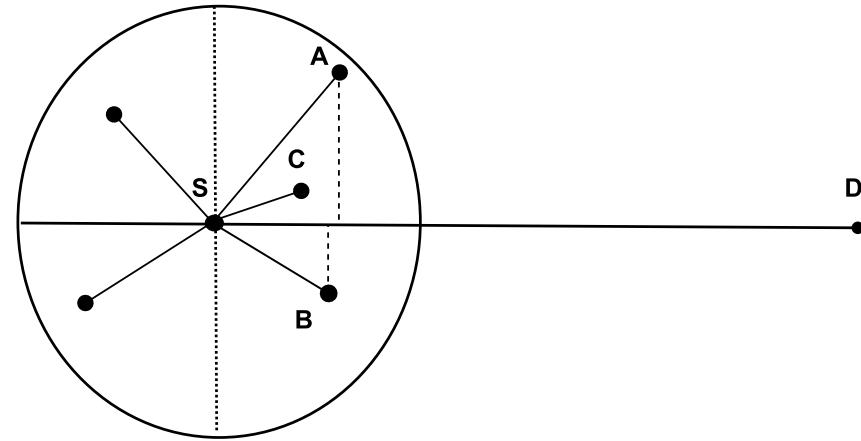
- MAC in early packet radio networks
 - Randomized access, no channel reservation
 - ALOHA, slotted ALOHA, CSMA
- Reservation-based protocols
 - Control packet exchange to reserve channel
 - MACA, MACAW, IEEE 802.11 DCF
- Power aware protocols
 - Power control, adjust power level to reach the receiver
 - Power management, allow nodes to turn off when idle, important in sensor networks
- Better performance by utilizing advanced hardware
 - Directional antennas
 - Multichannel transceivers

ROUTING

- Proactive protocols
 - Maintain an up-to-date network topology
 - Amount of routing traffic can be high
- Reactive protocols
 - Routes found and maintained on-demand
 - Less routing traffic but increased delay
- Hybrid protocols
 - Combine both proactive and reactive approaches
- Routing in sensor networks
 - Energy-efficiency, network lifetime maximization
 - Data-centric communications, the id of the original sender may be irrelevant
- Geographic routing

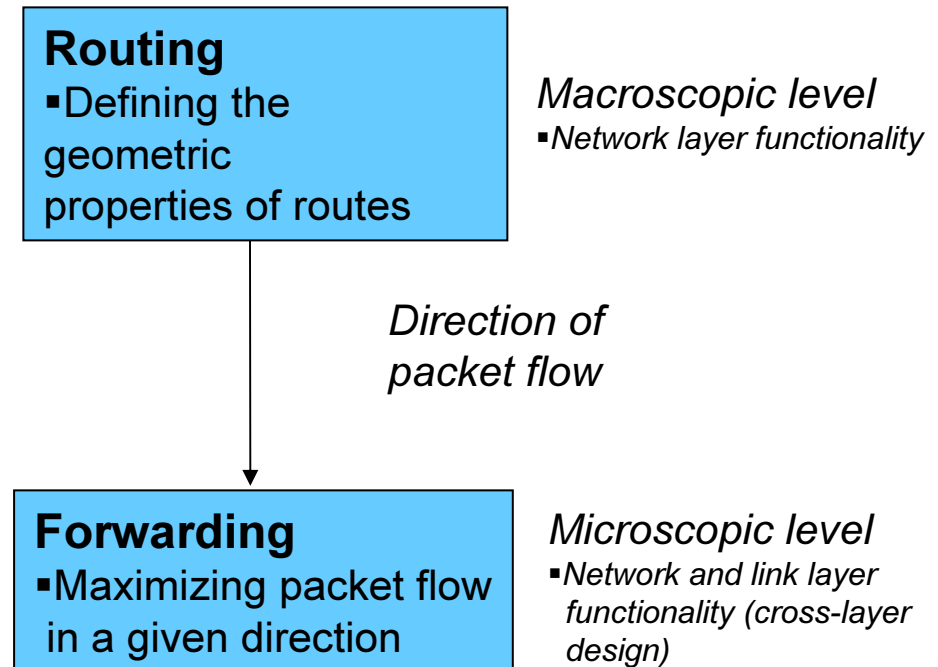
GEOGRAPHIC ROUTING

- Greedy forwarding
 - $\text{Progress}(A)$, $\text{distance}(B)$, $\text{angle}(C)$
- Routing around concave nodes
 - Face routing based on Gabriel graphs
- Location service
 - Responds to queries about the location of a destination



NETWORK MODEL

- Nodes distributed according to the two-dimensional Poisson point process with intensity λ
- Boolean interference model with a fixed transmission range R
 - Collision if a receiver hears more than one transmissions
- Slotted ALOHA MAC protocol with transmission probability p
- Each node knows its own and neighbors' location as well as the direction of packet flow



MEAN DENSITY OF PROGRESS

- The packet flow intensity is maximized

$$I = \rho v_x \quad [1/(\text{m}\cdot\text{s})],$$

where ρ is the packet density [$1/\text{m}^2$] and v_x is the mean packet velocity projected to the direction of the packet flow

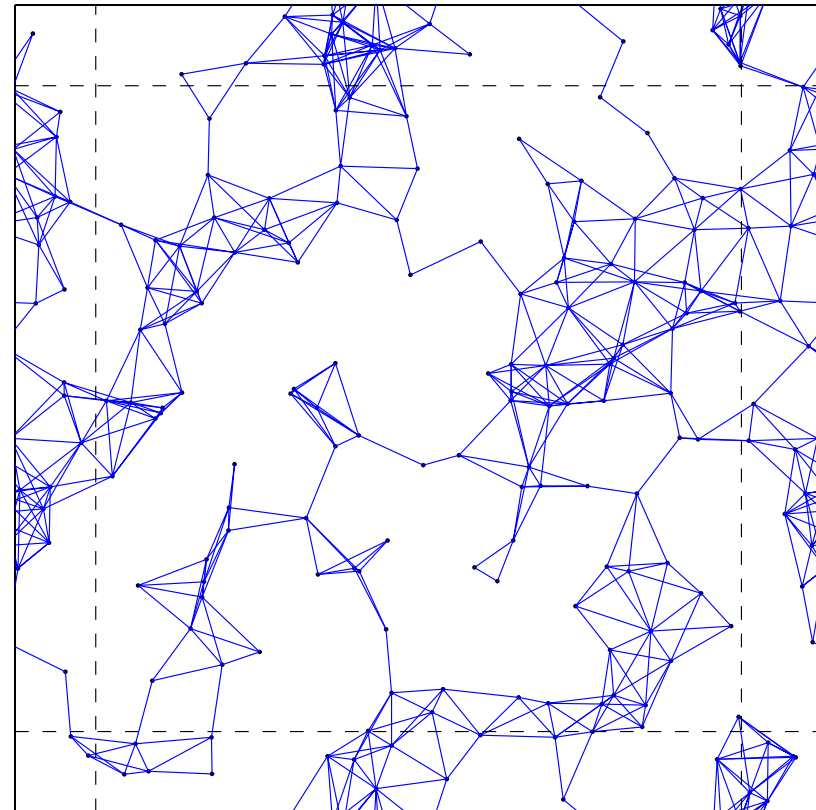
- Alternative definition for I : mean density of progress

$$I = (\sqrt{\lambda}/t) \cdot u(N_R, p) \quad [1/(\text{m}\cdot\text{s})],$$

where t is the time slot length [s], $N_R = \lambda\pi R^2$ is the average degree of a node and $u(N_R, p)$ is the mean progress of packets per time slot per node measured with $1/\sqrt{\lambda}$ as a unit length

SIMULATION MODEL

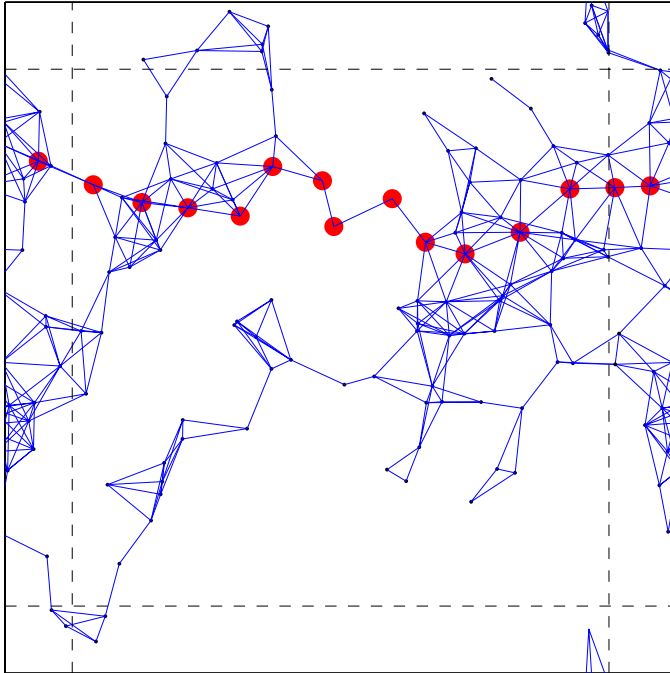
- Surface of the network plane seamed together into a torus
- Heavy traffic by initially placing 50 packets in each node
- Packets have infinite lifetime, no new packets generated
- Implementation using C++



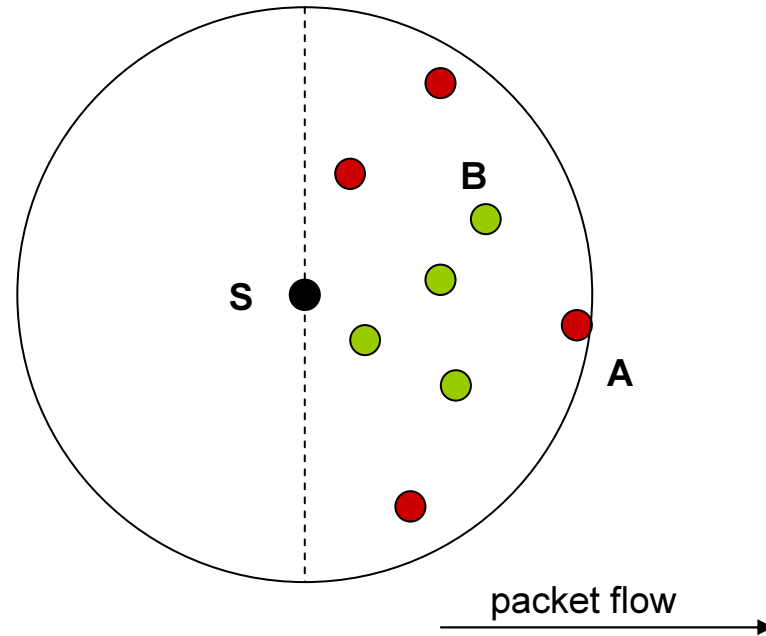
USED FORWARDING ALGORITHMS 1/2

- Most forward within radius (**MFR**)
 - Packet is forwarded to the most forward neighbor
- Random forwarding (**RF**)
 - Packet is forwarded to a random forward neighbor
- Weighted random forwarding (**WRF**)
 - Packet is forwarded to forward neighbor i with a probability q_i that is weighted with the progress from sender to i
- Opportunistic forwarding (**OF**)
 - Packet is forwarded to all forward neighbors
 - The most forward neighbor that successfully received the packet accepts the packet, others drop it

USED FORWARDING ALGORITHMS 2/2



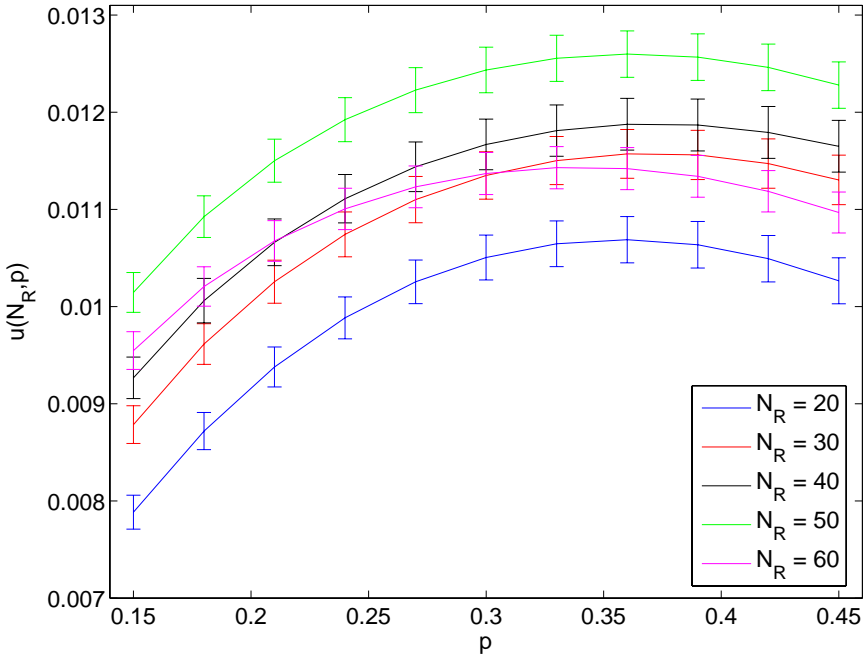
- MFR forwards packets into static paths



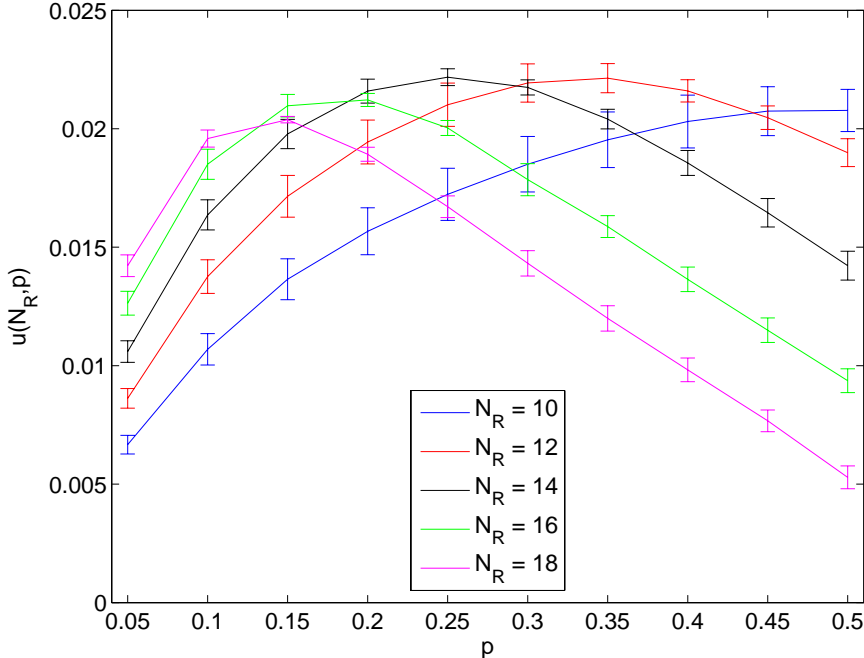
- MFR forwards to A, collision
- OF forwards to B, success
- RF may forward to any node, success probability 1/2

RESULTS 1/2

MFR

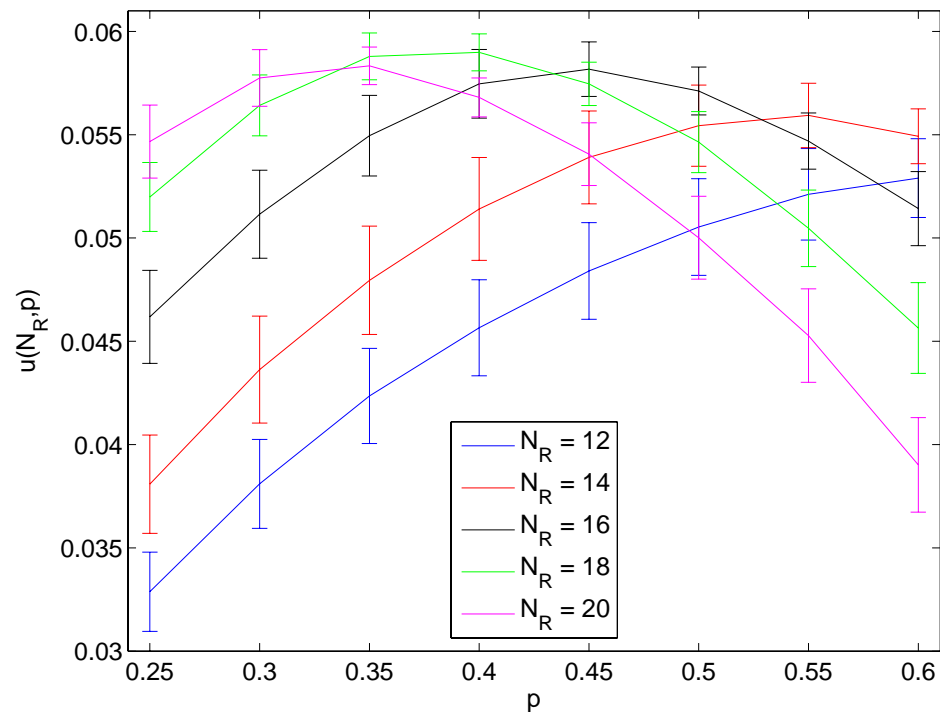


RF



RESULTS 2/2

OF

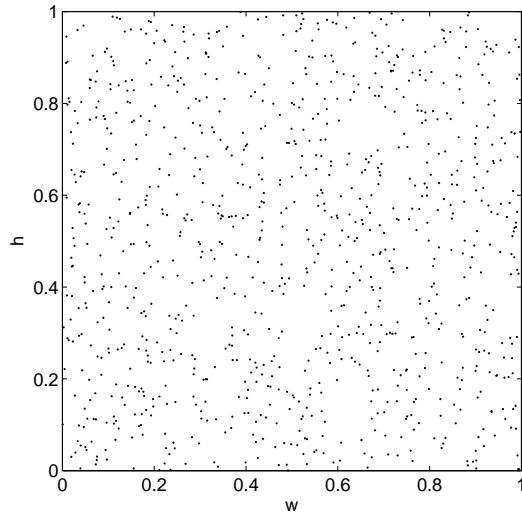


Maximum $u(N_R, p)$

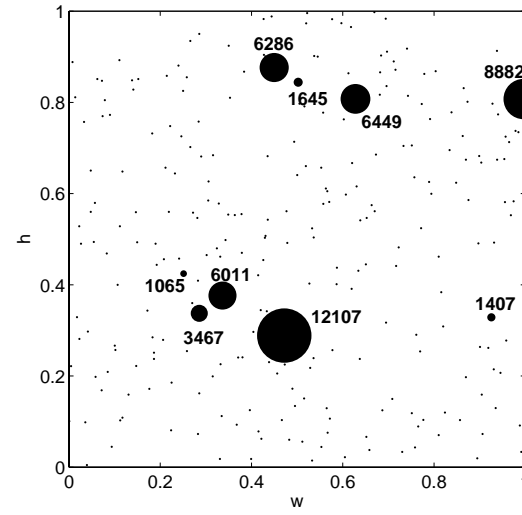
	$u(N_R, p)$	N_R	p
MFR	0.0126	50	0.35
RF	0.0222	14	0.25
WRF	0.0279	14	0.3
OF	0.059	18	0.4

- Opportunistic forwarding achieves clearly the best performance

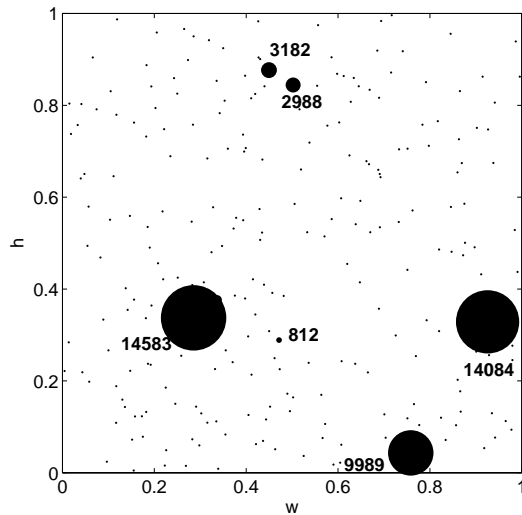
DISTRIBUTION OF PACKETS



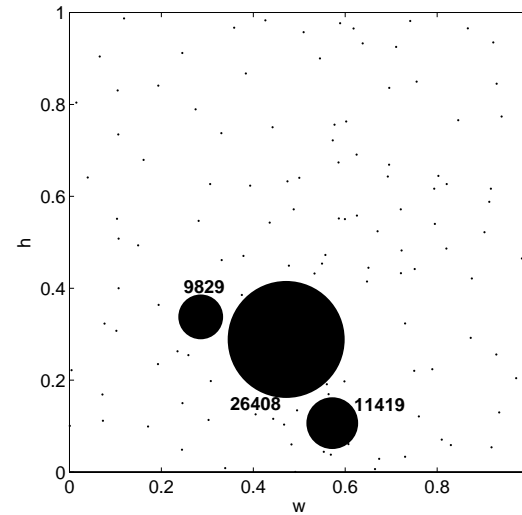
Node locations



RF



WRF



OF

CONCLUSIONS AND FURTHER WORK

- Randomized forwarding performs better than deterministic in a large ad hoc network
- Opportunistic forwarding improves throughput significantly
- Potential further work:
 - Take also into account the queue sizes at neighboring nodes when forwarding
 - Effect of power control
 - Effect of a more realistic interference model
 - Effect of node mobility